Badger: Fast Key-Value DB in Go

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Dgraph Labs

- Fast, Distributed graph database.
- Sparse data sets.
- Lots of relationships.

https://dgraph.io

What is Badger?

- Badger is an embedded key-value database, written in Go.
- Licensed under Apache 2.0.

go get github.com/dgraph-io/badger/...

Current Status

- Closing v2.0.
- Close to 3500 Github starts.
- 42 contributors.
- Used by Dgraph, Go-IPFS, 0-stor, Sandglass.

Serving 300TB (and growing) at Usenet Express

Basic Operations

Set a key-value

```
func set() error {
    fmt.Println("\nRunning SET")
    return db.Update(func(txn *badger.Txn) error {
        if err := txn.Set([]byte("foo"), []byte("bar")); err != nil {
            return err
        }
        fmt.Println("Set foo to bar")
        return nil
    })
}
```

Get a key-value

```
func get() error {
    fmt.Println("\nRunning GET")
    return db.View(func(txn *badger.Txn) error {
        item, err := txn.Get([]byte("foo")) // handle err
        if err != nil {
            return err
        }
        val, err := item.Value() // handle err
        if err != nil {
            return err
        }
        fmt.Printf("The value is: %s\n", val)
        return nil
    })
}
```

Iterate key-values

```
func iterate() error {
    fmt.Println("\nRunning ITERATE")
    return db.View(func(txn *badger.Txn) error {
        opts := badger.DefaultIteratorOptions
        it := txn.NewIterator(opts)
        defer it.Close()
        for it.Rewind(); it.Valid(); it.Next() {
            k := it.Item().Key()
            v, err := it.Item().Value() // handle err
            if err != nil {
                return err
            fmt.Printf("key=%s, value=%s\n", k, v)
        return nil
   })
```

Run the code

```
func main() {
    opt := badger.DefaultOptions
    opt.Dir = "/tmp/db"
    opt.ValueDir = opt.Dir
    var err error
    db, err = badger.Open(opt)
    if err != nil {
        panic(err)
    }
    defer db.Close()
    fmt.Println("DB opened")
    set()
    get()
    iterate()
    fmt.Println("DB done")
}
```

Badger != replacement for Go map

Motivation and Outcome

Cgo is not Go

Some people, when confronted with a problem, think "I know, I'll use cgo."

Now they have two problems.

-> Cgo is not Go, Dave Cheney

RocksDB

- Great write throughput.
- Okay read throughput.

Cons:

• Required Cgo.

BoltDB

- Pure Go.
- Great read throughput.

Cons:

• Bad write throughput.

Why build it?

- Go native key-value DB for Dgraph.
- No compromise in read-write performance.
- Avoid Cgo.

What did we spend?

- Spent a few months.
- Built with <1 full-time gopher.
- Aka, the power of Go!

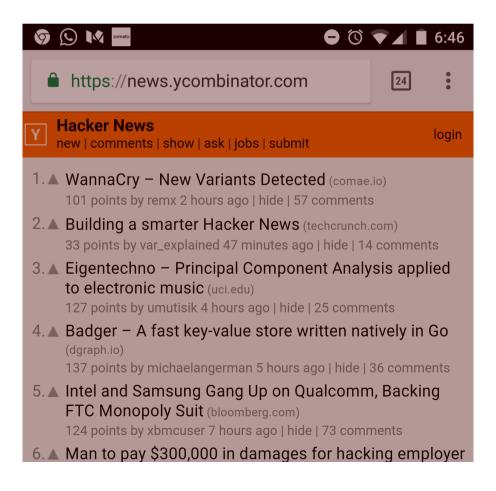


What did we gain?

- A faster key-value DB for Go.
- Ability to run Go profilers all the way down to disk.
- Clean Go code (no C).

Launch Reception

- Within 12 hours of blog post release
- First page of HN for a day.
- 355 points, 96 comments.
- 1250 Github stars in 4 days.



(bbc.com) 22 points by happy-go-lucky 2 hours ago | hide | 14 comments 7. A Why We Are Self-Publishing the Aviary Cookbook (medium.com) 31 points by wnm 2 hours ago | hide | 6 comments 8. Makers of Crowdfunded "Gravity Blanket" Withdraw Unsupported Medical Claims (statnews.com) 30 points by simon_acca 4 hours ago | hide | 26 comments 9. A Startup School 12: Alan Kay, Part II [video] (startupschool.org) 44 points by sama 4 hours ago | hide | 13 comments 10. A Mathematician's Lament (2002) [pdf] (maa.org) 89 points by Tomte 8 hours ago | hide | 16 comments 11. A Data Science Workflow: Overview and Challenges (acm.org)

Recruiters loved it!

• Got 3 different emails from 3 different recruiters...

Recruiters loved it!

• For jobs in the same company.

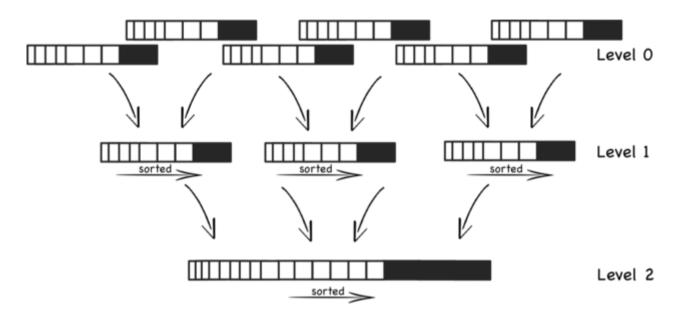


Design

Two common Trees

- LSM trees
- B+ trees

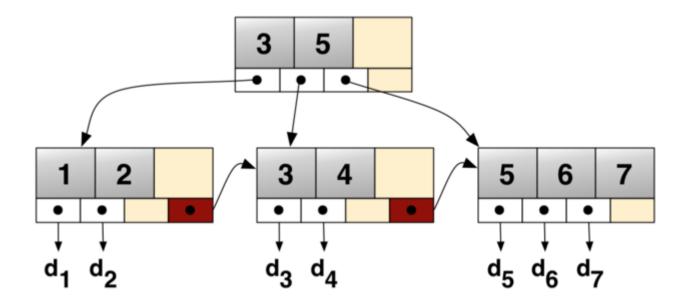
LSM Trees



Compaction continues creating fewer, larger and larger files

- More levels
- High write throughput
- High read latency
- Example: RocksDB

B+ Trees



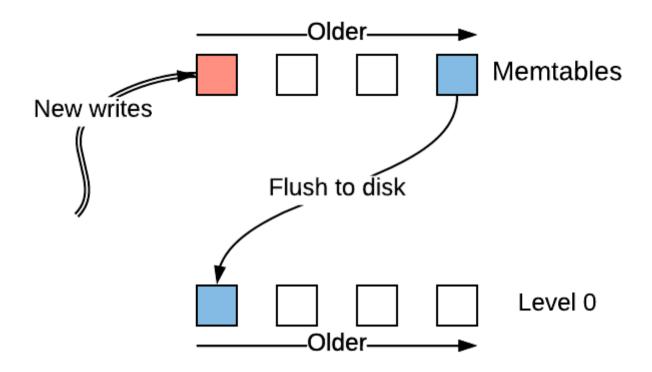
- Fewer levels
- Low write throughput
- Low read latency
- Example: BoltDB

Badger is based on LSM trees.

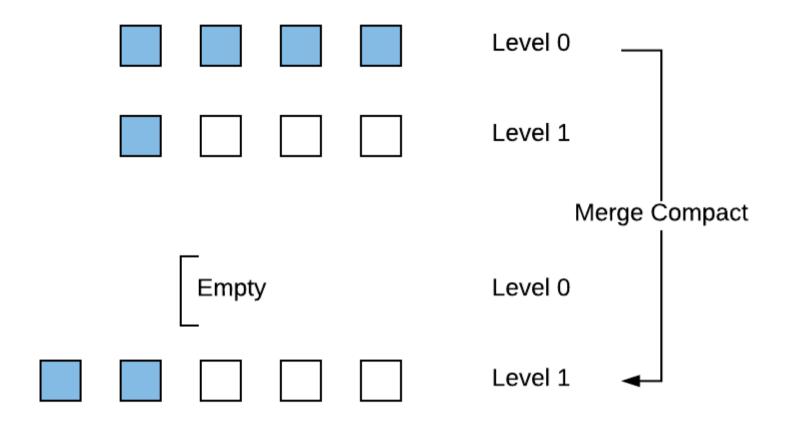
LSM Trees

Memtable (64 MB)				
L0 (256 MB)				
L1 (256 MB)				
L2 (2.5 GB)				10x
L3 (25 GB)				> previous level
L3 (250 GB)				

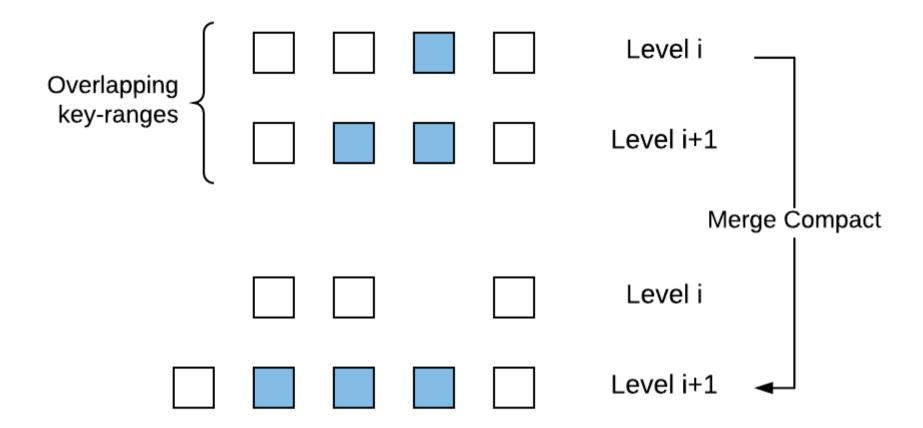
Writes in LSM trees: Memtable to L0



Writes in LSM trees: L0 to L1



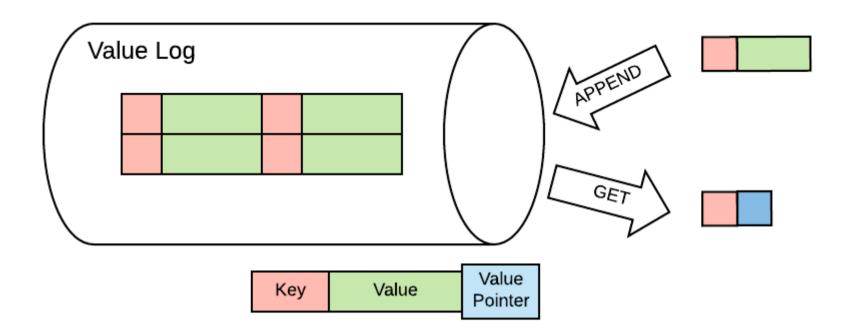
Writes in LSM trees: Li to Li+1



What makes Badger unique?

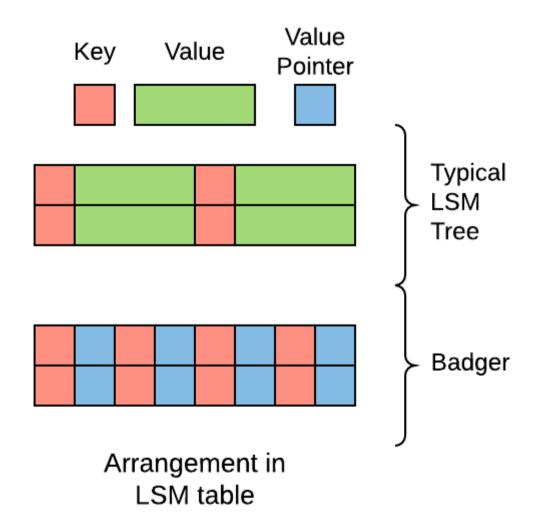
- Based on WiscKey paper by Uni Wisconsin-Madison.
- Separates keys from values.
- Stores keys in LSM tree.
- Stores value in value log.

Write to Value Log

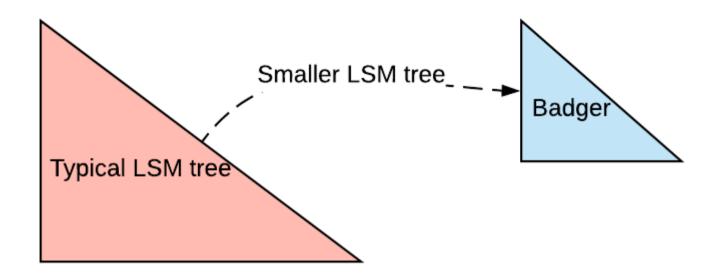


• Write value, get pointer.

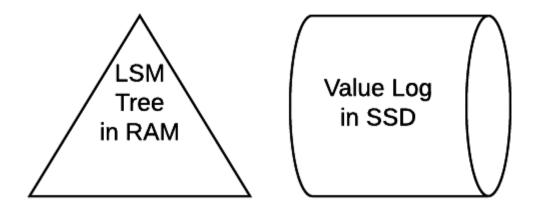
More keys per table



Smaller LSM tree



Typical Badger setup



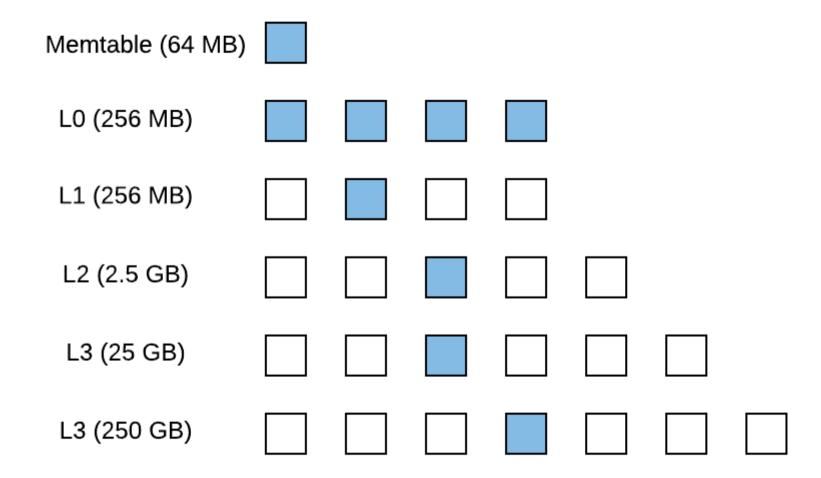
Advantages of smaller LSM tree

- Can be kept in RAM.
- Low read amplification (fewer lookups).
- Low write amplification (fewer compactions).
- ∝ Number of keys.

Usenet Express

- Hundreds of terabytes of data.
- Few gigabytes of LSM tree.

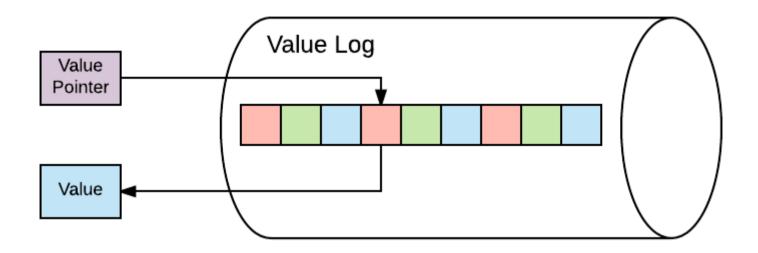
Reads in Badger: LSM tree



Reads in Badger: Bloom Filters



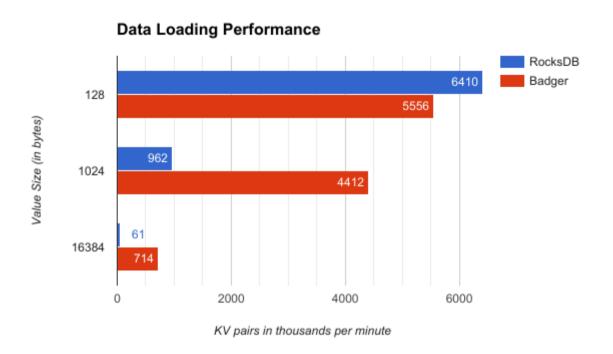
Reads in Badger: Value Log



• Once key found in LSM tree, read from value log.

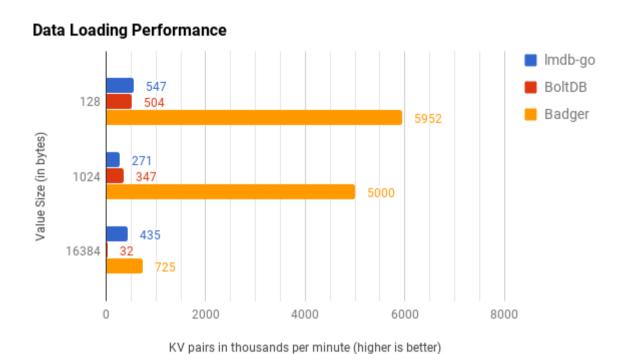
Badger is FAST-er

Data Loading: Badger vs Go-RocksDB



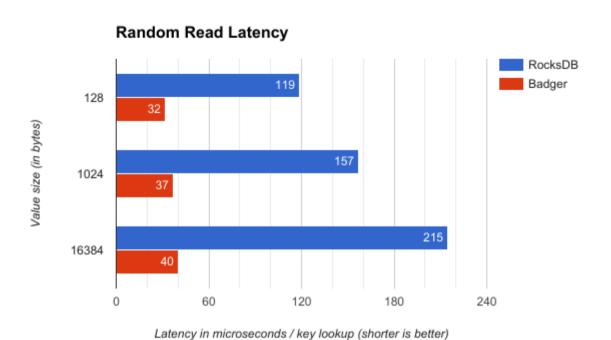
• As value size increases, Badger's becomes 11.7x faster.

Data Loading: Badger vs BoltDB



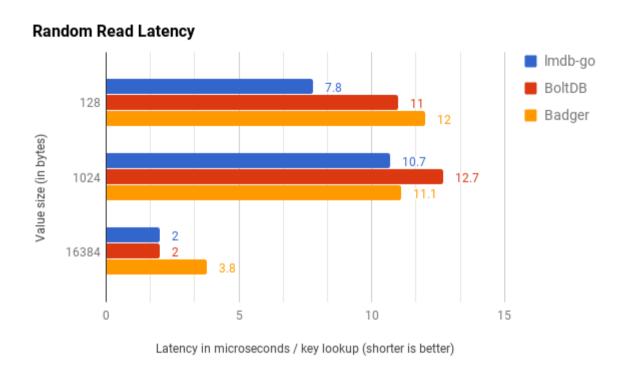
• 11x - 22x faster than BoltDB on all value sizes.

Random Reads: Badger vs Go-RocksDB



• Random Get latency is 3.7x - 5.3x lower than RocksDB.

Random Reads: Badger vs BoltDB



Random Get latency is slightly better or worse, depending on value size.

Various other benchmarks

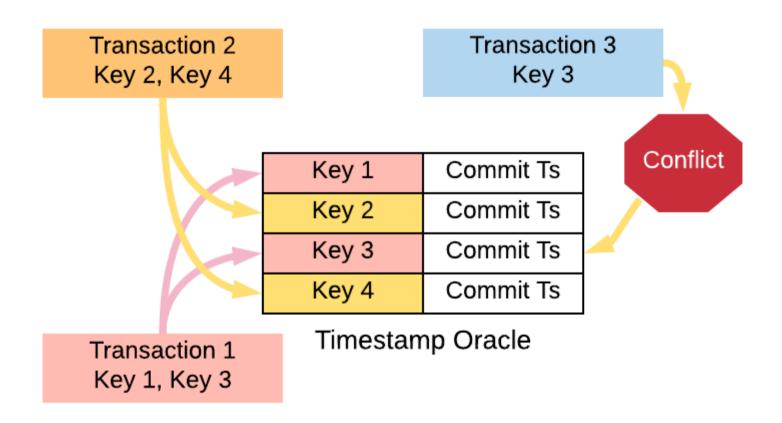
- Range iteration latency, etc.
- Can be found on https://blog.dgraph.io/
- Benchmarking code is open sourced.

github.com/dgraph-io/badger-bench (https://github.com/dgraph-io/badger-bench)

Features

Concurrent Transactions

• Badger uses Oracle to achieve concurrent lock-free transactions.



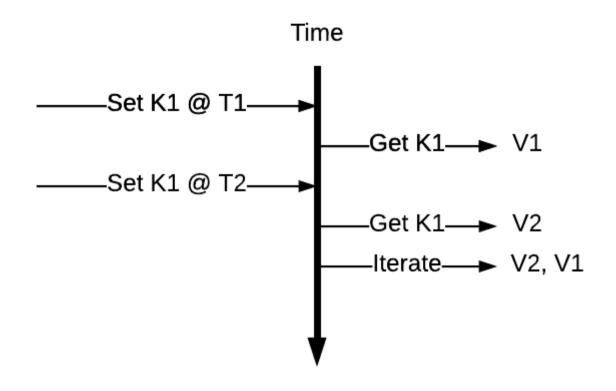
Concurrent Writes

- Batch up writes from multiple transactions.
- Amortize cost of disk write.
- No wait -> Smart Batching.

Smart Batching in Go

```
reqs := make([]*request, 0, 10)
                                          closedCase:
for {
                                              close(db.writeCh)
                                              for r := range db.writeCh {
    var r *request
    select {
                                                  reqs = append(reqs, r)
                                              }
    case r = <-db.writeCh:</pre>
    case <-lc.HasBeenClosed():</pre>
         goto closedCase
                                              pendingCh <- struct{}{}</pre>
                                              writeRequests(reqs)
                                              return
    for {
         reqs = append(reqs, r)
                                         writeCase:
         reqLen.Set(int64(len(reqs)))
                                              go writeRequests(reqs)
                                              reqs = make([]*request, 0, 10)
        select {
                                              reqLen.Set(0)
         case r = <-db.writeCh:</pre>
         case pendingCh <- struct{}{}:</pre>
             goto writeCase
        case <-lc.HasBeenClosed():</pre>
             goto closedCase
```

Multi Version Concurrency Control



- Badger stores multiple versions of the key.
- Provides direct access to the versions, via iterate.

Crash Resilience

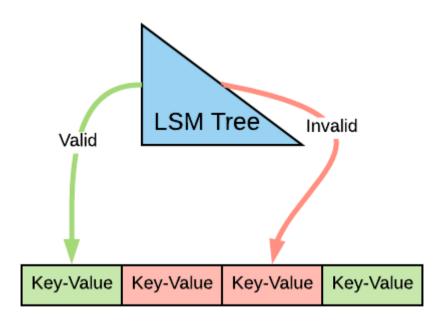
- LSM Memtables can be lost to crashes.
- Can be recovered from value log on restart.

Value Log Garbage Collection

Why?

- Value log would keep growing with every Set.
- Older versions of keys can be deleted.
- Corresponding values can be deleted from value log.

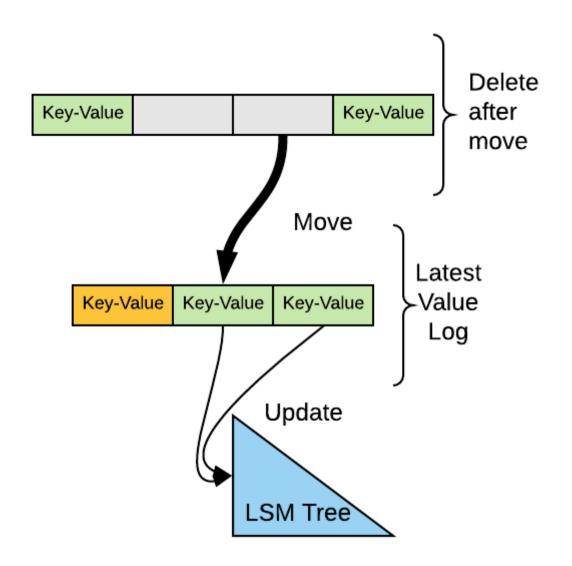
Stage 1: Punch Holes (v2.0 in Linux)



Punch Holes



Stage 2: Move entries, Delete log



Dealing with Qu-err-key file systems.

Would a file delete reclaim space in the file system?

Delete, no reclaim

No

```
if err := t.fd.Truncate(0); err != nil {
    // This is very important to let the FS know
    // that the file is deleted.
    return err
}
```

• Truncate the file before deleting.

Would closing a file sync its contents to disk?

Close, no-sync

No

```
if err := lf.fd.Sync(); err != nil {
    return errors.Wrapf(err, "Unable to sync value log: %q", lf.path)
}
if err := lf.fd.Close(); err != nil {
    return errors.Wrapf(err, "Unable to close value log: %q", lf.path)
}
```

• Explicitly sync file before closing.

Can a new synced file be lost?

Create, no-found

Yes

```
f, err := os.Open(dir)
if err != nil {
    return errors.Wrapf(err, "While opening directory: %s.", dir)
}
err = f.Sync()
closeErr := f.Close()
if err != nil {
    return errors.Wrapf(err, "While syncing directory: %s.", dir)
}
return errors.Wrapf(closeErr, "While closing directory: %s.", dir)
```

• Sync a directory just like you would sync a file.

Can a crash add garbage data to end of file?

Crash, no-clean

- Yes.
- Add checksums to know when to truncate a file.

Who should use Badger?

Don't use Badger if...

- You no Go! (C++, Java)
- You have a single-threaded sequential workload.
- You have a small, read-only workload.
- All your data can fit in memory easily.

Use Badger if...

- You Go!
- You want to avoid Cgo.
- You want a performant read-write workload.
- You access data concurrently (many goroutines accessing data).
- You need 3-dimensional access.

Future Work

- Encryption at rest.
- Others? (tell us what you need)

Work on Badger and Dgraph. Come join us!

github.com/dgraph-io/badger(https://github.com/dgraph-io/badger)

github.com/dgraph-io/dgraph(https://github.com/dgraph-io/dgraph)

Careers at Dgraph (https://dgraph.io/about.html)

Talk to us on Wechat



Thank you

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